Groups of automorphisms of rooted trees I

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Georg-August-Universität Göttingen

- 1. Introduction
- 2. Automorphisms of regular rooted trees
- 3. Branch groups
- 4. Next lecture

Introduction

- Milnor's Problem \implies growth of a group.
- General Burnside Problem \implies finiteness properties of a group.

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Grigorchuk's answer (1980):

Yes, the first ... Grigorchuk group.

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In modern terminology the general Burnside problem asks:

can a finitely generated periodic group be finite?

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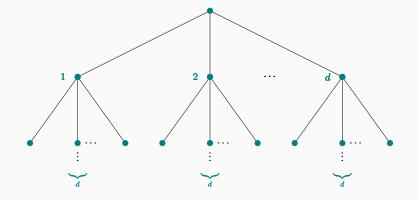
Milnor's question \bigcap General Burnside Problem

= the first Grigorchuk group, \ldots

Automorphisms of regular rooted trees



Seriously: the regular rooted tree \mathcal{T}_d



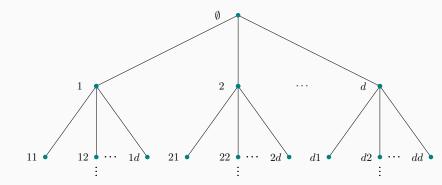
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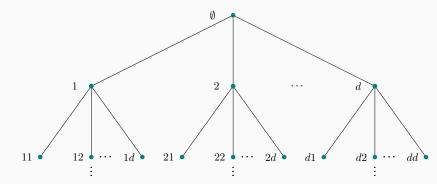
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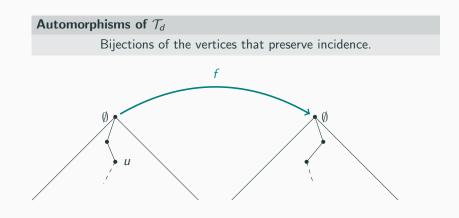
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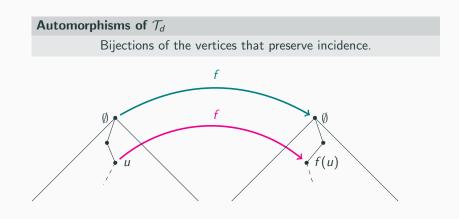


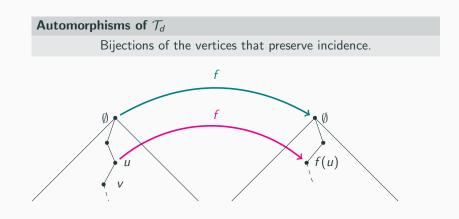
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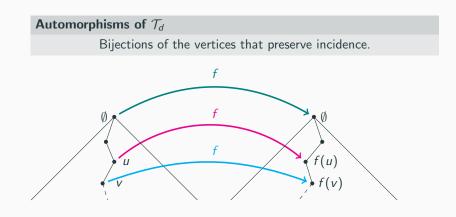


• X^n denotes the *n*th level of the tree.









The set Aut \mathcal{T}_d of all automorphisms of \mathcal{T}_d is a group with respect to composition between functions.

The set Aut T_d of all automorphisms of T_d is a group with respect to composition between functions.

Sometimes we write \mathcal{T} for \mathcal{T}_d , and, consequently, Aut \mathcal{T} for Aut \mathcal{T}_d .

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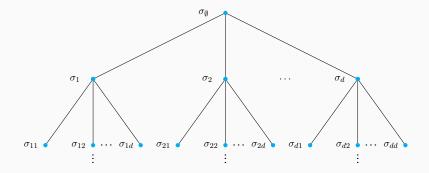
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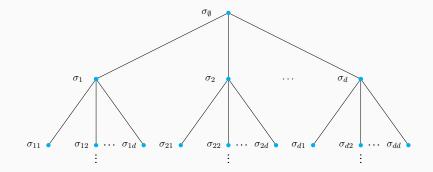
• Hence Aut T is a residually finite group (i.e. a group in which the intersection of all its normal subgroups of finite index is trivial).

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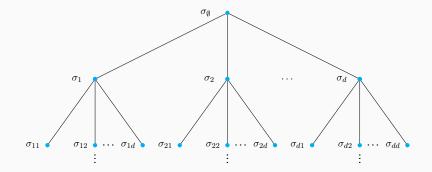


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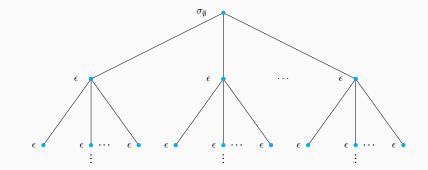
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We say that $\sigma_v \in \text{Sym}(d)$ is the *label* of f at the vertex v. The set of all labels is the *portrait* of f.

The simplest type are rooted automorphisms: given $\sigma \in \text{Sym}(d)$, they simply permute the *d* subtrees hanging from the root according to σ .

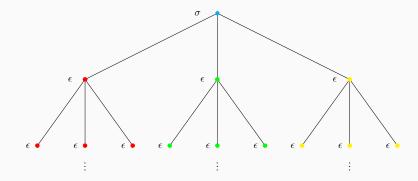


We denote with ϵ the identity element of Sym(*d*).

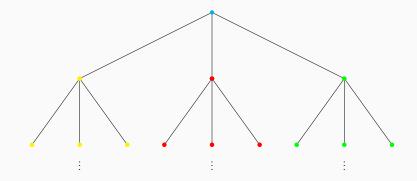
Let \mathcal{T}_3 be the ternary tree, and *a* the rooted automorphism corresponding to the cycle $\sigma = (1 \ 2 \ 3)$.

Example of a rooted automorphism

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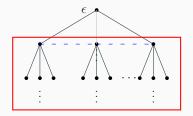
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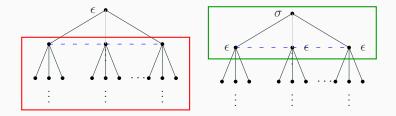
Note: sometimes we will identify a with σ .

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Some facts about $\operatorname{Aut} \mathcal{T}$: II

We define the isomorphism

$$\psi: \mathsf{st}(1) \longrightarrow \mathsf{Aut}\,\mathcal{T} imes \stackrel{d}{\cdots} imes \mathsf{Aut}\,\mathcal{T}$$
 $g \longmapsto (g_1, \dots, g_d)$

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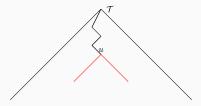
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Digression: this implies that Aut \mathcal{T} contains products Aut $\mathcal{T} \times \cdots \times Aut \mathcal{T}$. • Any $g \in \operatorname{Aut} \mathcal{T}_d$ can be seen as

 $g = h\sigma, \quad \sigma \in \operatorname{Sym}(d), \quad h \in \operatorname{st}(1) \cong \operatorname{Aut} \mathcal{T}_d \times \overset{d}{\ldots} \times \operatorname{Aut} \mathcal{T}_d$

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In other words, every $f \in \operatorname{Aut} \mathcal{T}_d$ can be written as

$$f=(f_1,\ldots,f_d)a,$$

where $f_i \in Aut \mathcal{T}_d$ and a is rooted corresponding to some permutation $\sigma \in Sym(d)$.

Example

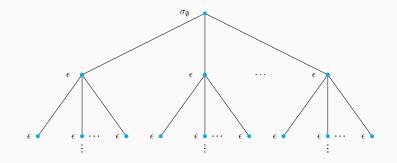
Let $f \in \text{Aut } \mathcal{T}_d$ with $f = (f_1, f_2, \dots, f_d)a$, where $f_i \in \text{Aut } \mathcal{T}_d$ and a is rooted corresponding to σ . If $f_1 = f_2 = \dots = f_d = 1$, then f is rooted.

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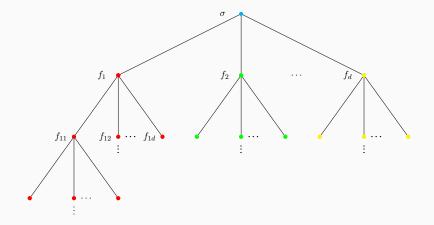
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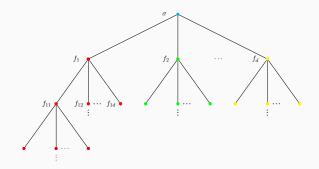
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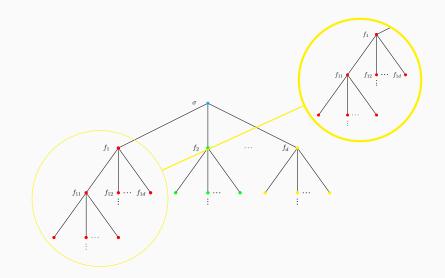


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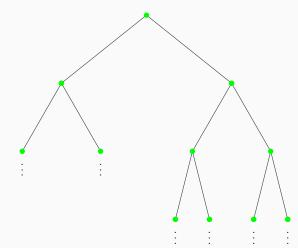


If \mathcal{T}_2 is the binary tree and ${\it a}$ is rooted corresponding to (1 2), let

b = (1, b)a.

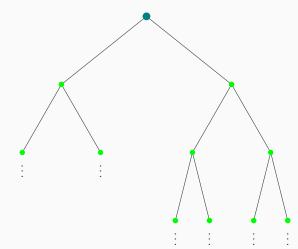
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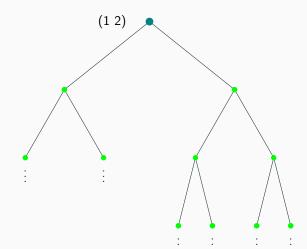
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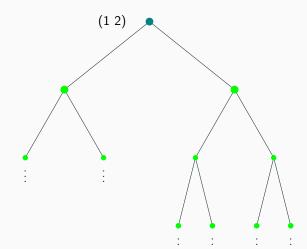
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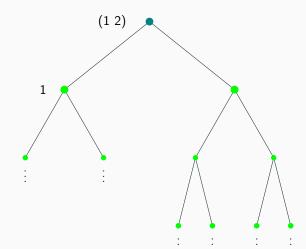
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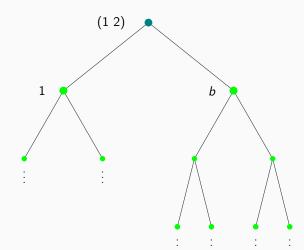
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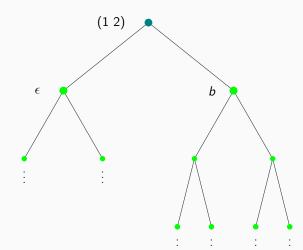
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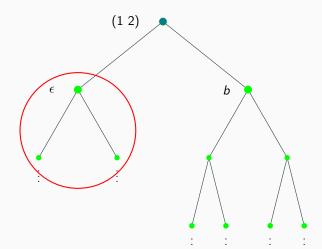
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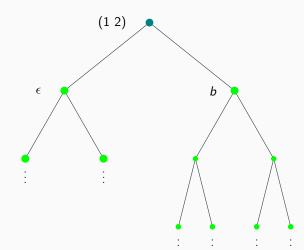
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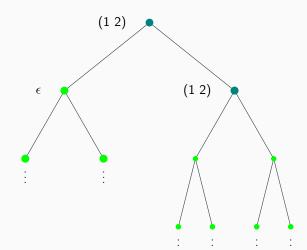
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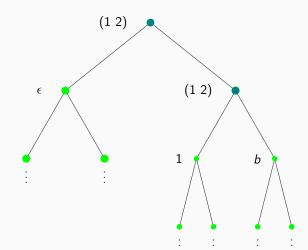
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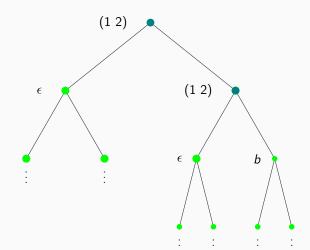
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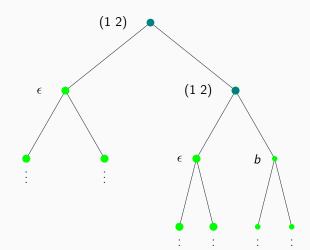
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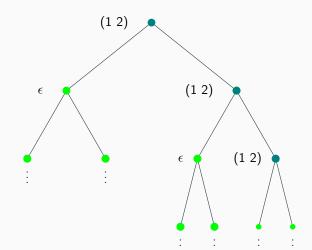
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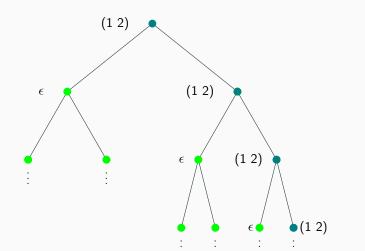
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How does *b* act on \mathcal{T}_2 ?



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$$\operatorname{st}(n) \simeq \operatorname{Aut} \mathcal{T} \times \stackrel{d^n}{\cdots} \times \operatorname{Aut} \mathcal{T},$$

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 The question is: given G ≤ Aut T, can we find for every n ∈ N a subgroup (eventually of finite index) of st_G(n) which is a direct product?

Rigid stabilizers

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 $\mathsf{rst}_G(u) = \{g \in G : g \text{ fixes all vertices outside } \mathcal{T}_u\}$

n-th level



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The *rigid stabilizer* of the *n*th level is $rst_G(n) = \prod_{u \in X^n} rst_G(u)$.

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- $\bullet\,$ Bad news: this is not usually the case for arbitrary subgroups of $\operatorname{Aut}\mathcal{T}.$
- Good news: in some cases, there exist "nice" rigid stabilizers.
- Informally speaking: the subgroup ψ_n(rst_G(n)) is the largest subgroup of ψ_n(st_G(n)) which is a "geometric" direct product.

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• We say that G is a branch group if for all $n \ge 1$, the index of the rigid *n*th level stabilizer in G is finite. In other words, for all $n \ge 1$,

 $|G: \operatorname{rst}_G(n)| < \infty.$

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- The first Grigorchuk group is a branch group.

Let $G \leq \operatorname{Aut} \mathcal{T}$.

 A group G is said to be *self-similar* if taken g = (g₁,...,g_d)σ ∈ G we have g_i ∈ G for any i = {1,...,d}. Let $G \leq \operatorname{Aut} \mathcal{T}$.

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- Example: Aut ${\mathcal T}$ is self-similar, the first Grigorchuk is self-similar.
- Non-example: The group G = ⟨a, b⟩, where a = (b, c)σ and c ∉ G, then G is not self-similar.

Regular branch groups

Let G be a self-similar group. We say that G is a *regular branch* if there exists a subgroup K of $st_G(1)$ of finite index such that

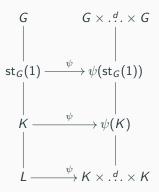
 $\psi(K) \supseteq K \times \stackrel{d}{\ldots} \times K.$

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More precisely we have this situation:



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- Regular branch \longrightarrow branch.

Next lecture

Next week we will present the following groups of automorphisms of rooted trees together with their main properties:

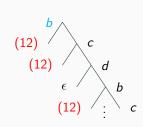
- The Grigorchuk groups
- The GGS-groups
- The Basilica group
- The Hanoi Tower group

I am sure I was too quick, so there is still time

 $\Gamma = \langle a, b, c, d \rangle$ $a = (1, 1)(12) \qquad b = (a, c) \qquad c = (a, d) \qquad d = (1, b)$

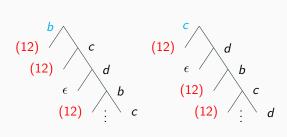
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$$a = (1,1)(12)$$
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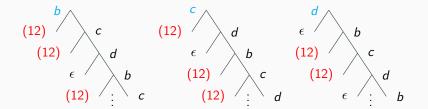
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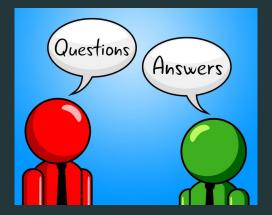
a = (1,1)(12) b = (a,c) c = (a,d) d = (1,b)



Some properties of $\ensuremath{\mathsf{\Gamma}}$

29th of October

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Thank you :)