

Equations in groups that are virtually direct products

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Abstract

In this note, we show that the satisfiability of equations and inequations with recognisable constraints is decidable in groups that are virtually direct products of finitely many hyperbolic groups.

Dedicated to Charles Sims, who introduced the first author to equations in groups.

2010 Mathematics Subject Classification: 20F10, 20F67, 68Q45

Key words: equations in groups, hyperbolic group, direct products, wreath products.

1 Introduction

For any group G and set of variables \mathcal{Y} , an *equation* with coefficients g_1, \dots, g_{m+1} from the group G is a formal expression $g_1 Y_1^{\epsilon_1} g_2 Y_2^{\epsilon_2} \dots Y_m^{\epsilon_m} g_{m+1} = 1_G$, where $\epsilon_j = \pm 1$ for all $1 \leq j \leq m$, and $Y_j \in \mathcal{Y}$. Such an equation is called *satisfiable* if there exist values for the Y_j 's in G with which the above identity in G is satisfied; each such set of values for the Y_j is a *solution*. Analogously, an *inequation* has the form $g_1 Y_1^{\epsilon_1} g_2 Y_2^{\epsilon_2} \dots Y_m^{\epsilon_m} g_{m+1} \neq 1_G$. A finite set of equations and inequations with coefficients in G is a *system of equations and inequations* over G , and is satisfiable if there are assignments to the Y_j such that all of the equations and inequations in the system are satisfied. For a group G , we say that systems of equations and inequations over G are *decidable* over G if there is an algorithm to determine whether any given such system is satisfiable. This question is widely known as the *Diophantine Problem* for G .

This article investigates equations in groups that are virtually direct products, and hence addresses the natural question of whether the decidability of equations extends from a group G to a group that contains G as a subgroup of finite index. While the decidability of equations in free groups was established in the 1980s by Makanin [18], it was only shown in 2010 that the same holds for virtually free groups: Dahmani and Guirardel [6] reduced the Diophantine Problem in virtually free groups to the same question relating to systems of twisted equations and inequations in free groups, and using difficult topological arguments they proved such systems to be decidable; a different approach to the Diophantine Problem in virtually free groups can be found in [17].

Here, in Theorem 3.1, we settle the Diophantine Problem in any group that is virtually a direct product of a finitely generated abelian group and non-elementary hyperbolic groups (equivalently, virtually a direct product of hyperbolic groups). We do this not by extending the result from the finite index direct product, but by embedding the whole

group into a direct product of permutational wreath products where the above questions can be answered positively (Lemma 3.5).

In fact, we prove Theorem 3.1 for an extended form of the Diophantine Problem, which asks if it can be decided whether a given system of equations and inequations has solutions in which some of the variables are constrained to lie in specified recognisable subsets of the group. (We define recognisable subsets in the next section.)

Our results show that the Diophantine Problem with recognisable constraints can be answered positively for, amongst others, dihedral (i.e. 2-generator) Artin groups, and groups that are virtually certain types of right-angled Artin groups. We note that, since any dihedral Artin group can alternatively be decomposed as a central extension of \mathbb{Z} by a virtually free group, decidability of its systems of equations (but not the more general problem with recognisable constraints) could also be derived from [16].

Acknowledgements

The authors acknowledge support from an LMS ‘Celebrating New Appointments’ grant for the meeting *Combinatorics and Computation in Groups*, where discussions on this paper started, and thank the ICMS for hosting the meeting. They would also like to thank Volker Diekert and Rick Thomas for clarifications and references on constraints.

2 Background and notation

Let G be a group with finite inverse closed generating set S , and let $\pi : S^* \rightarrow G$ be the natural homomorphism to G from the free monoid S^* generated by S . When w is a word over S , we write $|w|$ to denote the length of w as a word and $|\pi(w)|_G$ to denote the length of the shortest word over S that represents $\pi(w)$.

Definition 2.1.

- (1) A subset L of G is said to be recognisable if the full preimage $\pi^{-1}(L)$ is a regular subset of S^* .
- (2) A subset L of G is said to be rational if L is the image $\pi(L')$ of a regular subset L' of S^* .
- (3) A regular subset L' of S^* is quasi-isometrically embedded (*q.i. embedded*) in G if there exist $\lambda \geq 1$ and $\mu \geq 0$ such that, for any $w \in L'$, $|\pi(w)|_G \geq \frac{1}{\lambda}|w| - \mu$.
- (4) A rational subset L of G is quasi-isometrically embeddable (*q.i. embeddable*) in G if there exists a quasi-isometrically embedded regular subset L' of S^* such that $\pi(L') = L$.

It follows immediately from the definition that recognisable subsets of G are rational.

By [20, Proposition 6.3] a subset of G is recognisable if and only if it is a union of cosets of a subgroup of finite index in G , and hence a union of cosets of a normal subgroup of finite index (the core of a finite index subgroup will be both normal in G and of finite index in G); it follows that recognisability of a subset of G is independent of the choice of generating set for G . Rational subsets of G can be alternatively characterised as those sets that can be built out of finite subsets of G using finite union $A \cup B$, product $A \cdot B$, and semigroup generation A^* ; it follows from this that rationality of a subset of G is also independent of the choice of generating set.

By [1, Theorems 3.1, 3.2], a subgroup H of G is rational if and only if it is finitely generated, and recognisable if and only if it has finite index; the first of these results is attributed to Anisimov and Seifert.

We will be interested in the decidability of systems of equations and inequations in which some of the variables are restricted to lying in specified recognisable subsets of the group; this is the *Diophantine Problem with recognisable constraints*. This problem was considered for free groups and graph products in [8], [10] and [11].

Furthermore, if there exists an efficient algorithm that produces the solutions together with some concise and useful description of the solution set, we say that the system is *soluble* over G , or alternatively, that we can solve the *Diophantine Search Problem* in G . We observe that this second definition is intrinsically imprecise; in particular, for countable groups G with soluble word problem, we can in principle enumerate the solutions by the naive method of testing all m -tuples of elements of G .

The Diophantine Search Problem was solved in free groups by the work of Makanin and Razborov [19], and descriptions of the solutions are possible via Makanin-Razborov diagrams, or as EDT0L formal languages [3]. Then [7] shows the solubility of equations with rational constraints in virtually free groups, and [4] the solubility of equations with q.i. embeddable rational constraints in hyperbolic groups. Further, we observe that the set of solutions of a single equation or inequation over the free abelian group \mathbb{Z}^n generated by $X = \{x_i : 1 \leq i \leq n\}$ can be expressed as a deterministic context-free language over the alphabet $X \cup X^{-1}$. So the solution set of a system of equations and inequations over \mathbb{Z}^n is the intersection of finitely many such languages.

Our main result will rely on the following two statements about decidability of equations.

Proposition 2.2. (i) *Let G be a hyperbolic group. Then systems of equations and inequations with quasi-isometrically embeddable rational constraints are decidable in G .*

(ii) *Let G be a virtually abelian group. Then systems of equations and inequations with recognisable constraints are decidable in G .*

Proof. (i) This is proved in [6, Theorem 1].

(ii) We can adapt the proof of the same result, but without constraints, proved in [5, Lemma 5.4]. That proof reduces the problem in G to the same problem in a free abelian subgroup H of G . Given that any recognisable subset of G can be written as a union of cosets of a finite index subgroup M , we see that we can reduce our problem to the same problem in the free abelian subgroup $H \cap M$.

We note that systems of equations and inequations over the free abelian group \mathbb{Z}^n are decidable and soluble using standard techniques from integer linear algebra, such as the Hermite Normal Form for matrices. \square

Dihedral Artin groups. A dihedral Artin group DA_m , $m \geq 2$, is defined by the presentation

$$DA_m = \langle a, b \mid aba \cdots = bab \cdots \rangle,$$

where the single (braid) relation relates the two distinct alternating products of length m of the two generators a, b . When m is even, if we let $y_1 := a, y_2 := ab$, we see that

$$DA_m \cong \langle y_1, y_2 \mid y_1 y_2^{m/2} = y_2^{m/2} y_1 \rangle,$$

and hence we can describe DA_m as a central extension of the infinite cyclic group $\langle y_2^{m/2} \rangle$ by $\mathbb{Z} * C_{m/2}$. The latter group has $F_{m/2}$ as a subgroup of index $m/2$, and so in this case DA_m is virtually the direct product $\mathbb{Z} \times F_{m/2}$.

When m is odd, let $y_1 := ab \cdots a$, an alternating product of length m , and $y_2 := ab$. Then

$$DA_m \cong \langle y_1, y_2 \mid y_1^2 = y_2^m \rangle,$$

and so we can describe DA_m as a central extension of the infinite cyclic group $\langle y_2^m \rangle$ by $C_2 * C_m$; the free product has F_{m-1} as a subgroup of index $2m$. In this case, DA_m is virtually $\mathbb{Z} \times F_{m-1}$.

3 Main results

Theorem 3.1. *Let G be a finitely generated group that contains a direct product $A \times H_1 \times \cdots \times H_n$ as a finite index subgroup, where A is virtually abelian and H_1, \dots, H_n are non-elementary hyperbolic. Then systems of equations and inequations with recognisable constraints are decidable over G .*

We observe that since \mathbb{Z}^n is a direct product of elementary hyperbolic groups, G can also be expressed as virtually a direct product of hyperbolic groups. We make no assumption in this result that A is non-trivial or that n is non-zero.

Concerning solubility, it will be clear from the proof that descriptions of the solution sets over the factors H_i extend to descriptions in G .

Corollary 3.2. *Systems of equations and inequations with recognisable constraints are decidable in groups that are virtually dihedral Artin groups.*

Proof of corollary. As explained in Section 2, every dihedral Artin group is virtually a direct product of F_m and \mathbb{Z} , for $m \geq 2$. So the result follows from Theorem 3.1. \square

We note that we cannot expect to improve Theorem 3.1 to deal with rational constraints, as [9, Prop.11], due to Muscholl, shows. That result, concerning general right-angled Artin groups, implies in particular that systems of equations with rational constraints over direct products of non-abelian free groups are undecidable. The same holds for q.i. embeddable rational constraints, as Muscholl's constraints can be made to consist of geodesics.

We shall prove Theorem 3.1 at the end of the paper. It will follow from Proposition 2.2 and Theorem 3.3 below, once we have shown, in Proposition 4.4, that the group G of Theorem 3.1 contains an appropriate normal subgroup K .

For a group G , we define $\text{FIN}(G)$ to be the collection of groups each of which is either (isomorphic to) a subgroup of finite index in G or contains (a group isomorphic to) G as a subgroup of finite index. We introduce this terminology since we observe that, for G in various classes of groups that will be considered in this article, the decidability and solubility of systems of equations and inequations with various constraints holds not just in G but throughout $\text{FIN}(G)$.

Theorem 3.3. *Let G be a group and K a finite index normal subgroup, where $K = K_1 \times \cdots \times K_n$, and $\{K_1, \dots, K_n\}$ is a union of conjugacy classes of subgroups in G . Suppose that systems of equations and inequations with recognisable constraints are decidable in*

all groups in $\text{FIN}(K_i)$, for each i . Then systems of equations and inequations with recognisable constraints are decidable in G .

We need to establish some lemmas before proving this result.

Lemma 3.4. *Let $G = G_1 \times \cdots \times G_n$ be a direct product of groups G_i over which systems of equations and inequations with recognisable constraints are decidable. Then the same is true in G .*

Proof. Suppose that the system consists of a set E of equations and a set I of inequations. An equation from E has a solution in G if and only if, in each of the direct factors G_i , the projection onto G_i has a solution, while an inequation from I has a solution in G if and only, in at least one of the direct factors G_i , the projection onto G_i has a solution. Hence the system has a solution in G if and only if we can write I as a (not necessarily disjoint) union $I_1 \cup I_2 \cdots \cup I_n$, where for each i the projection of the system $E \cup I_i$ onto G_i has a solution in G_i . So decidability in G is inherited from decidability in the direct factors G_i .

Now suppose that some of the variables are restricted to lie in some specified recognisable subsets of G . We recall that each such subset is a finite union of cosets of a finite index normal subgroup M_j of G and by letting $H_i := \cap_j M_j \cap G_i$, we have $H_i \triangleleft G_i$, $|G_i : H_i| < \infty$, and $H := H_1 \times \cdots \times H_n$ is contained in all of the subgroups that arise in the constraints.

We first find all solutions of the projection of the system onto the finite quotient G/H of G . Then, for each such solution, the set of solutions of the original system that project onto it can be defined in terms of the solutions of a system of equations and inequations over G that are constrained to lie in H . As in the first paragraph of this proof, we can reduce the decidability of such a systems to the decidability of a finite collection of systems of equations and inequations over the component groups G_i for which the solutions are constrained to lie in H_i , and we can decide each of those by hypothesis. \square

The following lemma shows how one can embed an extension of a direct product into a direct product of wreath products, and has the flavour of a Kaloujnine-Krasner result [15] about embeddings of group extensions into wreath products; however, we do not specifically need that result here.

Lemma 3.5. *Let $K = K_1 \times \cdots \times K_n$ be a normal subgroup of finite index in a group G , and suppose that the set of subgroups $\{K_1, \dots, K_n\}$ is a union of conjugacy classes of subgroups in G .*

- (i) *If the subgroups K_i form a single conjugacy class, then G is isomorphic to a subgroup of finite index in $J \wr P$, where $J \cong N_G(K_1)/(K_2 \times \cdots \times K_n)$ contains a subgroup of finite index isomorphic to K_1 , and $P \leq \mathcal{S}_n$ is the image of the permutation action on $\{K_1, \dots, K_n\}$ induced by conjugation in G .*
- (ii) *Suppose that K_1, \dots, K_k are representatives of the conjugacy classes of K_1, \dots, K_n within G . Then G embeds as a subgroup of finite index in a direct product $W_1 \times W_2 \times \cdots \times W_k$ of wreath products $W_j = J_j \wr P_j$, where J_j is a group containing K_j as a subgroup of finite index and P_j is a finite permutation group.*

Proof. Part (i) is a rewording of [14, Theorem 4.1 (1)].

For (ii), for $1 \leq i \leq k$, let N_i be the product of those K_j that are not conjugate to K_i in G . Then $N_i \trianglelefteq G$, and we define $Q_i := G/N_i$ and let $\mu_i : G \rightarrow Q_i$ be the natural

map. Then the images under μ_i of those K_j that are conjugate in G to K_i form a single conjugacy class of subgroups of Q_i , and their product has finite index in Q_i . So by (i) Q_i embeds via a map η_i as a subgroup of finite index in a group $W_i = J_i \wr P_i$, where J_i contains $\mu_i(K_i) \cong K_i$ as a subgroup of finite index, and P_i is a finite permutation group.

Now the map $\mu : G \rightarrow W_1 \times \cdots \times W_k$ defined by $\mu(g) = (\eta_1\mu_1(g), \dots, \eta_k\mu_k(g))$ is an embedding of G into $W_1 \times \cdots \times W_k$. Since, for $i = 1, \dots, k$, $\mu(K_i)$ is a subgroup of the direct factor W_i isomorphic to K_i , and $\mu(K_i^G)$ has finite index in W_i , we see that $\mu(K)$ (and hence also $\mu(G)$) has finite index in $W_1 \times \cdots \times W_k$. \square

Lemma 3.6. *Suppose that systems of equations and inequations with recognisable constraints are decidable over the group J . Then they are also decidable over the permutation wreath product $W = J \wr P$ of J with a finite subgroup P of \mathcal{S}_n .*

Proof. Decomposing W as the split extension of n copies of J by $P \subset \mathcal{S}_n$, we represent each of its elements by an $(n+1)$ -tuple (j_1, \dots, j_n, π) , with $j_i \in J$, $\pi \in P$, with multiplication defined by

$$(j_1, \dots, j_n, \pi)(k_1, \dots, k_n, \rho) = (j_1 k_{\pi^{-1}(1)}, \dots, j_n k_{\pi^{-1}(n)}, \pi\rho).$$

For a given system of equations and inequations over W , we project this system onto the finite group P and find all of the finitely many solutions in P . For each such solution in P , we can use the displayed equation to reduce the problem of deciding whether there are solutions in W that project onto that particular solution in P to one of deciding a system of equations and inequations over the direct product J^n of n copies of J . We can do that by Lemma 3.4.

If the system over W has rational constraints, then this technique reduces the problem to one of deciding systems of equations with rational constraints over J^n , which we can again do by Lemma 3.4. \square

Proof of Theorem 3.3. Suppose that the groups K_1, \dots, K_n fall into k conjugacy classes under the conjugation action of G , of which K_1, \dots, K_k are representatives, and of the n original subgroups, n_j of them are conjugate to K_j , for each $1 \leq j \leq k$. Then, by Lemma 3.5, G embeds as a subgroup of finite index in a direct product $W_1 \times W_2 \times \cdots \times W_k$ of permutation wreath products $W_j = J_j \wr P_j$, where J_j is a group containing K_j as a subgroup of finite index and P_j is a subgroup of \mathcal{S}_{n_j} . Our hypotheses ensure that, for each $j = 1, \dots, k$, equations with recognisable constraints are decidable in J_j , and Lemma 3.6 ensures that the same is true in W_j .

By Lemma 3.4 systems of equations with recognisable constraints are decidable if that holds for the direct factors in the direct product $W_1 \times W_2 \times \cdots \times W_k$. Then since finite index subgroups are recognisable, it follows that the same is true within the finite index subgroup G of $W_1 \times \cdots \times W_k$. \square

4 Direct products of finite index

Recall that virtually cyclic groups (including finite groups) are hyperbolic, and are known as *elementary hyperbolic*. The following lemma lists the known properties of hyperbolic groups that we shall need.

Lemma 4.1. *Let H be a hyperbolic group. Then the centralizer of any element of H is hyperbolic, and is elementary if the element is non-torsion. Any subgroup of H consisting of torsion elements is finite, and there is a bound on the order of the finite subgroups of H . Furthermore, if H is non-elementary, then $Z(H)$ is finite and $H/Z(H)$ is non-elementary hyperbolic.*

Proof. The first two assertions are proved in [12, Proposition 4.3 and Proposition 5.1]. The finiteness of torsion subgroups is proved in [13, Corollaire 36, Chapitre 8] and the bound on the order of finite subgroups is proved in [2, Theorem III Γ .3.2]. It follows that $Z(H)$ is a torsion subgroup when H is non-elementary, so $Z(H)$ is finite, and the proof that $H/Z(H)$ is hyperbolic and non-elementary is straightforward. \square

Lemma 4.2. *Let A be virtually abelian, and H_1, \dots, H_n non-elementary hyperbolic groups. Let $H = A \times H_1 \times \dots \times H_n$, let $L \leq H$ with $|H : L|$ finite, and let $g \in L$. Suppose that the projection of g onto H_i is a non-torsion element of H_i for exactly k values of $i \in \{1, 2, \dots, n\}$. Then the centraliser of g in L is a subgroup of finite index in a group $B \times K_1 \times \dots \times K_j$, where B is virtually abelian, each K_i is non-elementary hyperbolic, and $j \leq n - k$.*

Proof. This follows from the previous lemma. \square

Lemma 4.3. *Let A and B be virtually abelian, and $H_1, \dots, H_m, K_1, \dots, K_n$ non-elementary hyperbolic groups. Let $H = A \times H_1 \times \dots \times H_m$ and $K = B \times K_1 \times \dots \times K_n$. Suppose that a group L is isomorphic to finite index subgroups of both H and K . Then $m = n$.*

Proof. Suppose that $m \leq n$ and use induction on m . If $m = 0$, then H is virtually abelian, and hence K must be virtually abelian, and so $n = 0$. So suppose that $m > 0$.

It is convenient to identify L with the subgroups of H and K with which it is isomorphic. By replacing H , K and L by finite index subgroups, we can assume that A and B are both abelian, and that L projects onto all of the direct factors H_i and K_i . Then $Z(L) = L \cap Z(H) = L \cap Z(K)$ and, since $L \cap A \leq Z(L)$ and $L \cap B \leq Z(L)$, it follows from Lemma 4.1 that $\bar{L} := L/Z(L)$ can be identified with finite index subgroups of $\bar{H}_1 \times \dots \times \bar{H}_m$ and of $\bar{K}_1 \times \dots \times \bar{K}_n$, where $\bar{H}_i := H_i/Z(H_i)$ and $\bar{K}_i := K_i/Z(K_i)$ are non-elementary hyperbolic groups.

Since $\bar{L} \cap \bar{K}_1$ has finite index in \bar{K}_1 it contains a non-torsion element g , and $C_{\bar{L}}(g)$ has finite index in a direct product of a virtually abelian group and $n - 1$ non-elementary hyperbolic groups. But, by considering g as an element of $\bar{H}_1 \times \dots \times \bar{H}_m$, we see from Lemma 4.2 that $C_{\bar{L}}(g)$ has finite index in the direct product of a virtually abelian group and $m - t$ non-elementary hyperbolic groups for some $t \geq 1$. So, by the inductive hypothesis, we have $m - t = n - 1$, and since $m \leq n$, we must have $t = 1$ and $m = n$. \square

Proposition 4.4. *Let A be virtually abelian, and H_1, \dots, H_n non-elementary hyperbolic groups. Suppose that a group G has a subgroup H of finite index with $H \cong A \times H_1 \times \dots \times H_n$. Then there is a normal subgroup K of finite index in G with $K \leq H$, such that $K \cong B \times K_1 \times \dots \times K_n$, where $B \leq A$ of finite index in A , each K_i is isomorphic to a finite index subgroup of H_i , and the set $\{B, K_1, \dots, K_n\}$ is a union of orbits under conjugation by G .*

Proof. By replacing H by a subgroup of finite index, we may assume that A is free abelian. Let N be the core of H in G ; then $N \trianglelefteq G$ has finite index and $N \leq H$.

Let $C := N \cap A$; then $|A : C| \leq |G : N| < \infty$. Also, $C \leq Z(N)$, and C is torsion free. The projection of $Z(N)$ onto each of the subgroups H_i is central in a subgroup of finite index in H_i and hence finite, so $|Z(N) : C| < \infty$. Let $k := |Z(N) : C|$ and define $B := Z(N)^k$. Then $B \leq C$ and B is characteristic and of finite index in $Z(N)$. So B has finite index in C , and hence in A , and since $Z(N)$ is normal in G , so is B .

The rest of this proof is devoted to the construction of the subgroups K_1, \dots, K_n of G . We find these as subgroups of finite index in subgroups L_1, \dots, L_n of G , which we identify by considering a quotient G/T of G , and considering its action by conjugation on its free abelian normal subgroup N/T .

We note that $[N, N] \leq H_1 \times \dots \times H_n$, so $B \cap [N, N] = 1$. Now choose T with $[N, N] \leq T \leq N$ so that $T/[N, N]$ is the torsion subgroup of the abelian group $N/[N, N]$; then N/T is free abelian. Since $T/[N, N]$ is characteristic in $N/[N, N]$, and $N/[N, N]$ is normal in G/N , certainly $T \trianglelefteq G$. Since B is torsion-free with $B \cap [N, N] = 1$, while $T/[N, N]$ has torsion, we have $T \cap B = 1$, and so $BT/T \cong B$. Furthermore, since B is normal in G , the image $BT/T \cong B$ of B in G/T is normal in G/T .

Let $g \mapsto \bar{g}$ denote the natural map from G to G/T . Then the conjugation action of \bar{G} on the free abelian group \bar{N} makes \bar{N} into a torsion-free $\mathbb{Z}\bar{G}$ -module in which \bar{BT} is a submodule. So by Lemma 4.5 below, there is a subgroup \bar{L} of \bar{N} with $\bar{L} \trianglelefteq \bar{G}$, $|\bar{N} : \bar{LBT}| < \infty$, and $\bar{L} \cap \bar{BT} = \{1\}$. Then, where $L \trianglelefteq G$ is the preimage of \bar{L} , we deduce that $|N : BL| < \infty$ (and so also $|H : AL| < \infty$) and $L \cap BT = T$. Since $T \cap B = 1$, we have $L \cap B = 1$, and also $L \cap A = 1$ (since if $g \in L \cap A$, we have $g^k \in L \cap B = \{1\}$, so $g = 1$, since A is torsion-free). Now the natural map from H to H/A , whose image is isomorphic to $H_1 \times \dots \times H_n$, maps AL to a group of finite index in H/A , which is isomorphic to L ; the image lifts to a subgroup M of finite index in $H_1 \times \dots \times H_n$, to which we associate an isomorphism $\phi : L \rightarrow M$. For each i , we define $L_i := \phi^{-1}(M \cap H_i)$. Then L_i is isomorphic to a subgroup of finite index in H_i , $L_1 \times \dots \times L_n$ has finite index in L , and $B \times L_1 \times \dots \times L_n$ has finite index in G .

Let h be a non-torsion element of L_i for some i . Then $C_L(h) \cong C_M(\phi(h))$ has finite index in the direct product of a virtually cyclic group and $n - 1$ non-elementary hyperbolic groups. Now, for any $g \in G$, the same applies to $h^g = g^{-1}hg$ and so, by Lemmas 4.2 and 4.3 applied to $\phi(h^g)$, we see that the projection of $\phi(h^g)$ onto H_j is a non-torsion element for exactly one value of j .

Furthermore, if h' is another non-torsion element of the same L_i , then $C_L(\langle h, h' \rangle)$ has finite index in the direct product of a (possibly finite) virtually cyclic group and $n - 1$ non-elementary hyperbolic groups, and again the same applies to $C_L(\langle h^g, h'^g \rangle)$. It follows that the unique subgroup H_j onto which the projection of $\phi(h^g)$ is a non-torsion element is the same as that onto which the projection of $\phi(h'^g)$ is a non-torsion element; hence we may denote that subgroup by H_{ig} . We notice too that we can find an integer r for which the projections of $\phi((h^g)^r) = \phi((h^r)^g)$ onto all subgroups H_j apart from H_{ig} are trivial. So we have $h^r \in L_i$, and $(h^r)^g \in L_{ig}$. Of course, for all g' , we also have $\phi(((h^r)^{gg'})) = \phi(((h^r)^g)^{g'})$. It follows that the unique subgroup H_r onto which the projection of $\phi(h^r)^{gg'}$ is a non-torsion element is identified both as $H_{igg'}$ and as $H_{(ig)^{g'}}$, and so we see that, for all g, g' we have $i^{gg'} = (ig)^{g'}$.

Now, for any element $h \in L_i$, whether or not it is a torsion element, and for any $j \neq ig$, the projection of $\phi(h^g)$ onto H_j is a torsion element. So the projection of $\phi(L_i^g)$ onto H_j is a torsion group and hence, by Lemma 4.1 is finite. Now let P_i denote the intersection of the kernels of all homomorphisms $\theta : L_i \rightarrow H_i$ for which $|\theta(L_i)| < \infty$. If $h \in P_i$, then $\phi(h^g) \in H_{ig}$ and hence $h^g \in L_{ig}$. Also, since by Lemma 4.1 there is a bound on the

orders of finite subgroups of H_i , $|L_i : P_i|$ is finite.

Finally, let $K_i = \{h \in L_i \mid h^g \in L_{ig} \forall g \in G\}$. Then it is straightforward to check that K_i is a subgroup of G and, since $P_i \leq K_i$, we see that $|L_i : K_i|$ is finite; hence K_i is isomorphic to a subgroup of finite index in H_i . It follows from the statement $i^{gg'} = (i^g)^{g'}$ that $K_i^g \leq K_{ig}$ for all i and g and then, since $(K_{ig})^{g^{-1}} \leq K_i$, we must have $K_i^g = K_{ig}$. So we have proved that $\{K_1, \dots, K_n\}$ is a union of orbits under the conjugation action of G .

This completes the proof. \square

Lemma 4.5. *Let G be a finite group, let V be a finite dimensional torsion-free $\mathbb{Z}G$ -module, and W a submodule. Then there exists a $\mathbb{Z}G$ -submodule U of V with $U \cap W = \{0\}$ such that $V/(U \oplus W)$ is finite.*

Proof. Let $\widehat{V} = V \otimes \mathbb{Q}$ and $\widehat{W} = W \otimes \mathbb{Q}$ be the corresponding $\mathbb{Q}G$ -modules. By Maschke's theorem, there exists a $\mathbb{Q}G$ -submodule \widehat{U} of \widehat{V} with $\widehat{V} = \widehat{U} \oplus \widehat{W}$. Let e_1, \dots, e_n be a \mathbb{Z} -basis of V , which we may consider also as a \mathbb{Q} -basis of \widehat{V} . We can choose a basis u_1, \dots, u_k of \widehat{U} such that the matrices representing the action of G have integer entries. Define $\lambda_{ij} \in \mathbb{Q}$ by $u_i = \sum_{j=1}^n \lambda_{ij} e_j$. Let m be a common multiple of the denominators of all the λ_{ij} , and define $U \subseteq V$ to be the \mathbb{Z} -module generated by the elements μ_i, \dots, μ_k of \widehat{U} . Then $U \oplus W$ has rank n , and so must have finite index in V . \square

Proof of Theorem 3.1. Let G be as in the hypothesis of the theorem. Then, by Proposition 4.4, G has a normal subgroup K of finite index, such that $K \cong B \times K_1 \times \dots \times K_n$, where $B \leq A$ with $|A : B|$ finite, each K_i is isomorphic to a finite index subgroup of H_i , and the set $\{B, K_1, \dots, K_n\}$ is a union of orbits under conjugation by G . Now all groups in $\text{FIN}(B)$ are virtually abelian, and all groups in $\text{FIN}(K_i)$ are hyperbolic for each i , so systems of equations and inequations with recognisable constraints are decidable in all groups in either $\text{FIN}(B)$ or $\text{FIN}(K_i)$ (any i). The result now follows by Theorem 3.3. \square

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